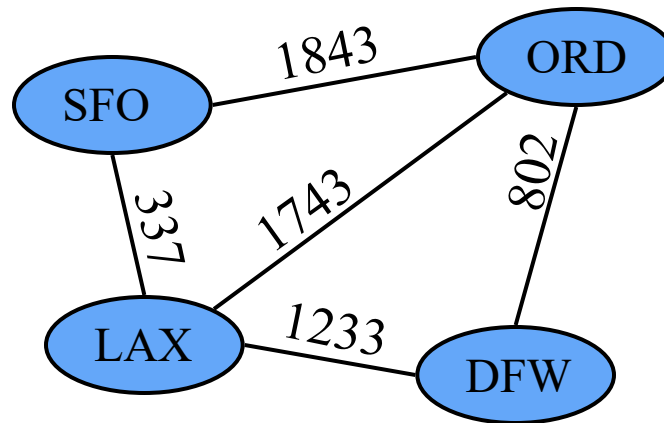


# Graphs



# Outline / Reading

## Graphs (6.1)

- Definition
- Applications
- Terminology
- Properties
- ADT

## Data structures for graphs (6.2)

- Edge list structure
- Adjacency list structure
- Adjacency matrix structure

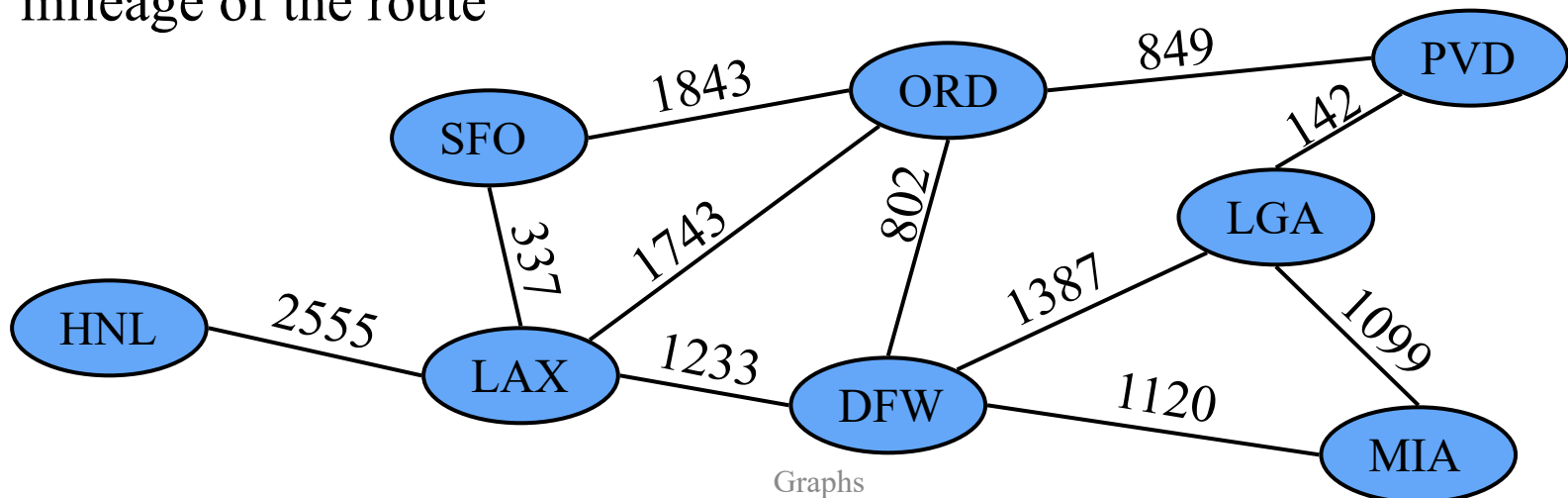
# Graph

A **graph** is a pair  $(V, E)$ , where

- $V$  is a set of nodes, called **vertices**
- $E$  is a collection of pairs of vertices, called **edges**
- Vertices and edges are positions and store elements

Example:

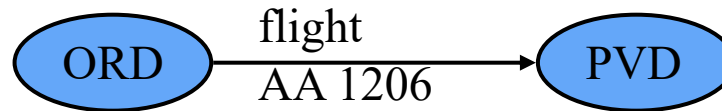
- A vertex represents an airport and stores the three-letter airport code
- An edge represents a flight route between two airports and stores the mileage of the route



# Edge Types

## Directed edge

- ordered pair of vertices  $(u, v)$
- first vertex  $u$  is the origin
- second vertex  $v$  is the destination
- e.g., a flight

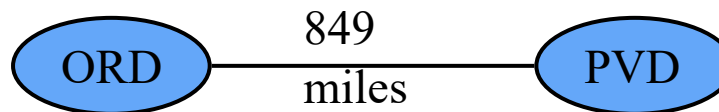


## Directed graph

- all the edges are directed
- e.g., flight network

## Undirected edge

- unordered pair of vertices  $(u, v)$
- e.g., a flight route

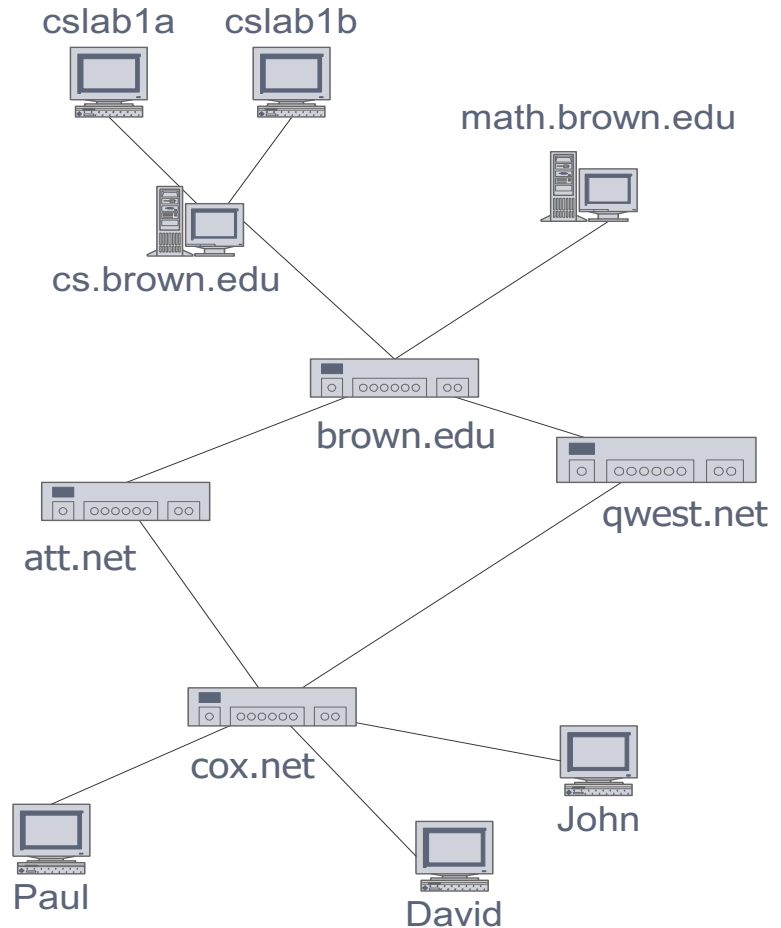


## Undirected graph

- all the edges are undirected
- e.g., route network

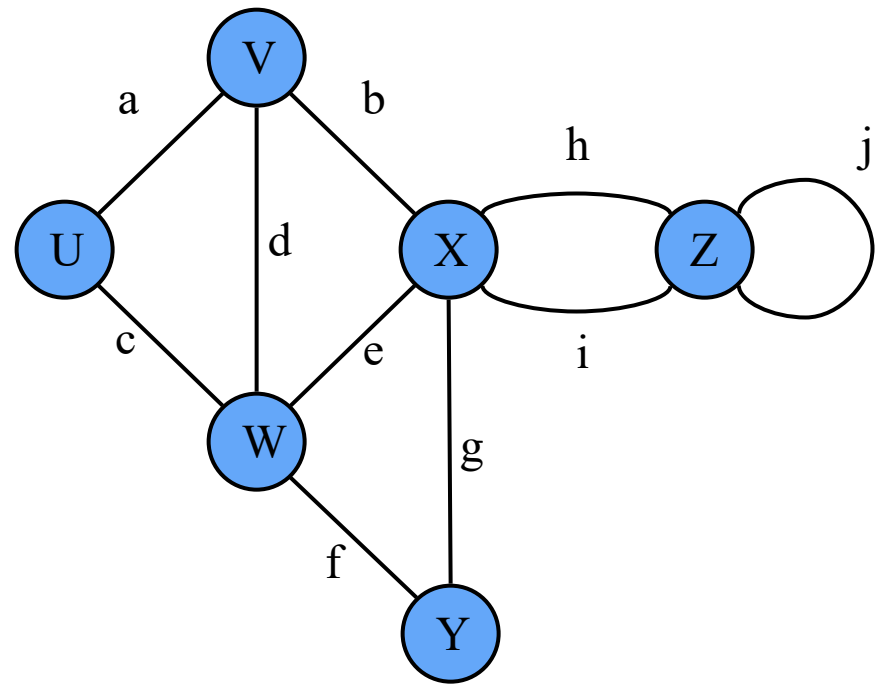
# Applications

- Electronic circuits
  - Printed circuit board
  - Integrated circuit
- Transportation networks
  - Highway network
  - Flight network
- Computer networks
  - Local area network
  - Internet
  - Web
- Databases
  - Entity-relationship diagram



# Terminology

- **End** vertices (or endpoints) of an edge
  - $U$  and  $V$  are the endpoints of  $a$
- Edges **incident** on a vertex
  - $a$ ,  $d$ , and  $b$  are incident on  $V$
- **Adjacent** vertices
  - $U$  and  $V$  are adjacent
- **Degree** of a vertex
  - $X$  has degree 5
- **Parallel** edges
  - $h$  and  $i$  are parallel edges
- **Self-loop**
  - $j$  is a self-loop



# Terminology (cont.)

## Path

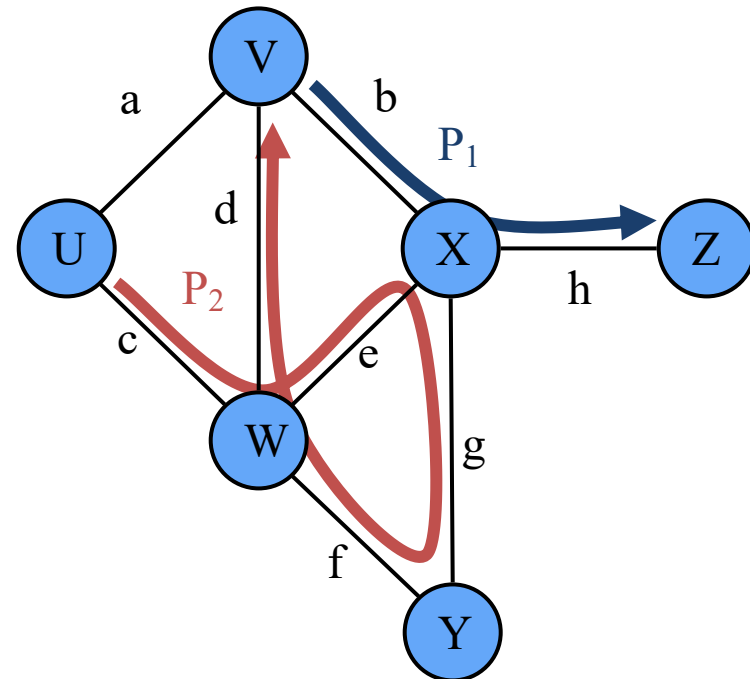
- sequence of alternating vertices and edges
- begins with a vertex
- ends with a vertex
- each edge is preceded and followed by its endpoints

## Simple path

- path such that all its vertices and edges are distinct

## Examples

- $P_1=(V,b,X,h,Z)$  is a simple path
- $P_2=(U,c,W,e,X,g,Y,f,W,d,V)$  is a path that is not simple



# Terminology (cont.)

## Cycle

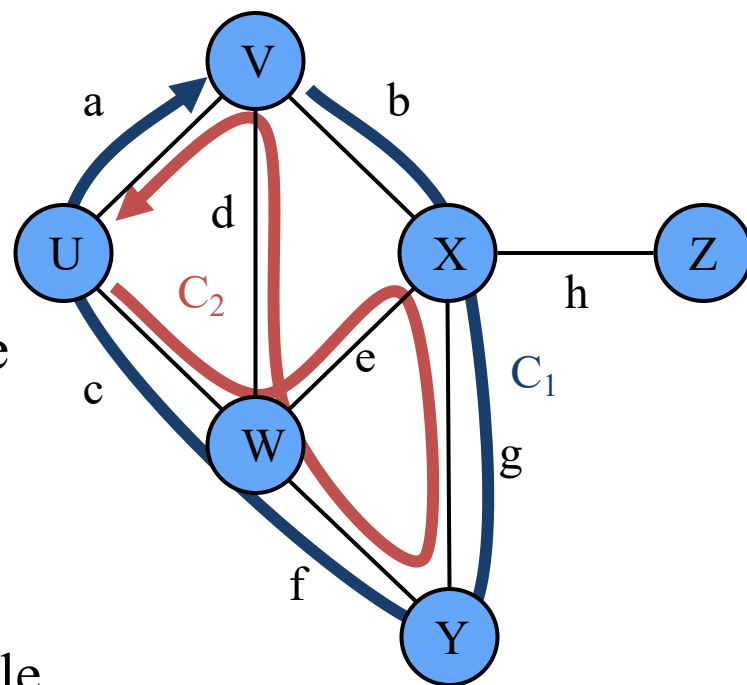
- circular sequence of alternating vertices and edges
- each edge is preceded and followed by its endpoints

## Simple cycle

- cycle such that all its vertices and edges are distinct

## Examples

- $C_1 = (V, b, X, g, Y, f, W, c, U, a, \downarrow)$  is a simple cycle
- $C_2 = (U, c, W, e, X, g, Y, f, W, d, V, a, \downarrow)$  is a cycle that is not simple





# Properties

**Property 1.** In an undirected graph

$$\sum_v \deg(v) = 2m$$

Proof: each edge is counted twice

**Property 2.** In an undirected graph with no self-loops and no multiple edges

$$m \leq n(n-1)/2$$

Proof: each vertex has degree at most  $(n-1)$

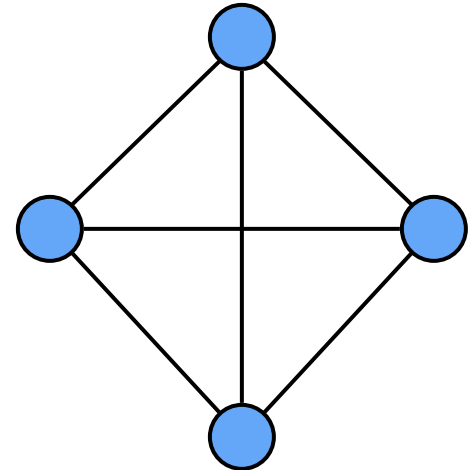
What is the bound for a directed graph?

## Notation

$n$  number of vertices

$m$  number of edges

$\deg(v)$  degree of vertex  $v$



Ex:  $n = 4$ ;  $m = 6$ ;  
 $\deg(v) = 3$

# Main Methods of the Graph ADT

## Vertices and edges

- are positions
- store elements

## Accessor methods

- `aVertex()`
- `incidentEdges(v)`
- `endVertices(e)`
- `isDirected(e)`
- `origin(e)`
- `destination(e)`
- `opposite(v, e)`
- `areAdjacent(v, w)`

## Update methods

- `insertVertex(o)`
- `insertEdge(v, w, o)`
- `insertDirectedEdge(v, w, o)`
- `removeVertex(v)`
- `removeEdge(e)`

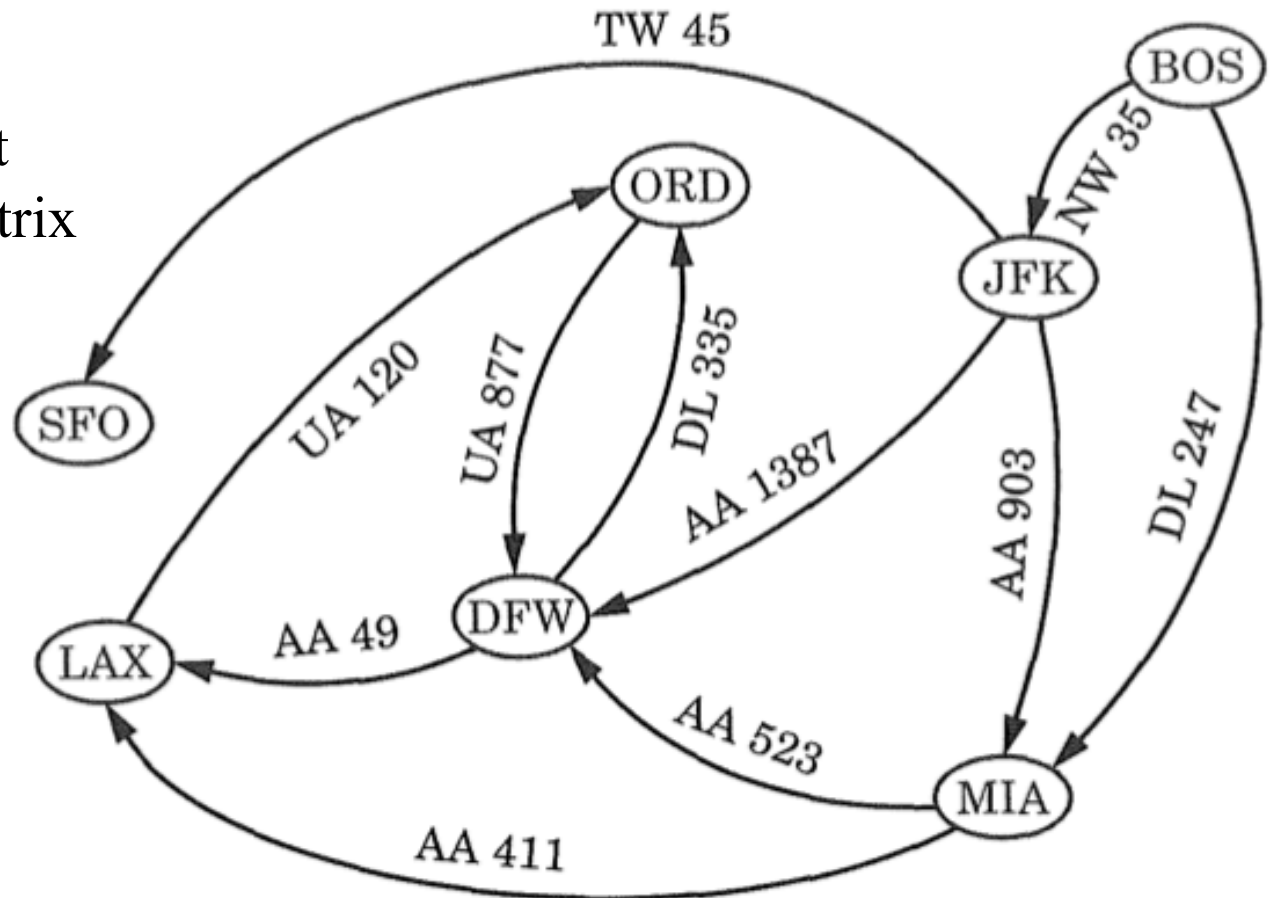
## Generic methods

- `numVertices()`
- `numEdges()`
- `vertices()`
- `edges()`

# Data Structures

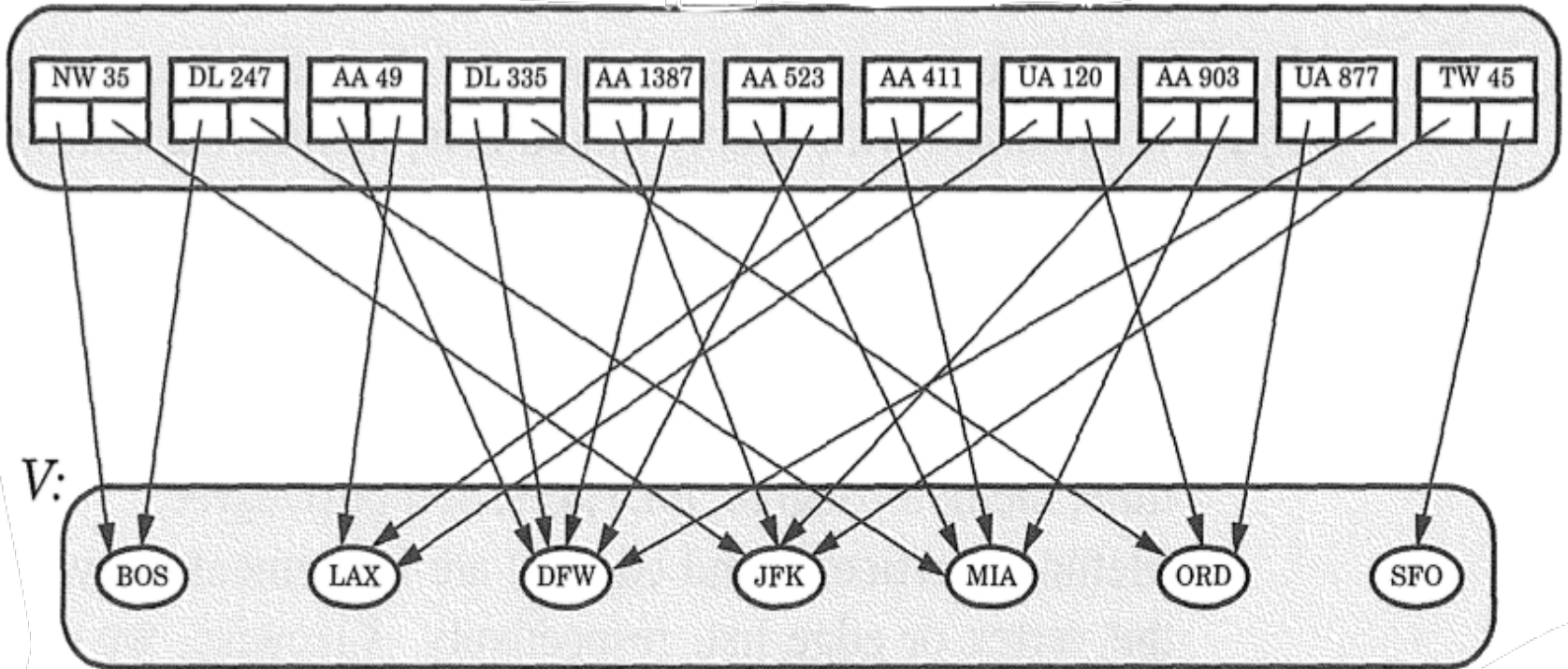
Structures to represent a graph:

1. Edge List
2. Adjacency List
3. Adjacency Matrix



# Edge List Structure

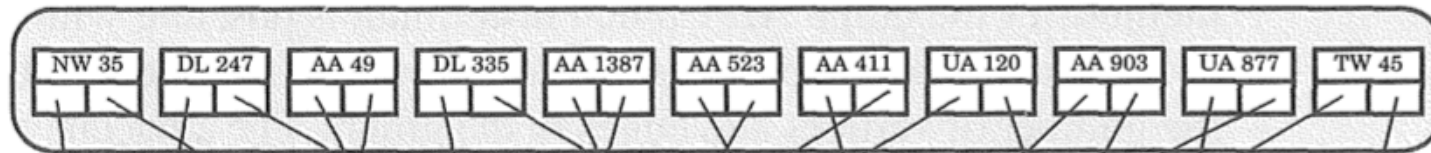
*E:*



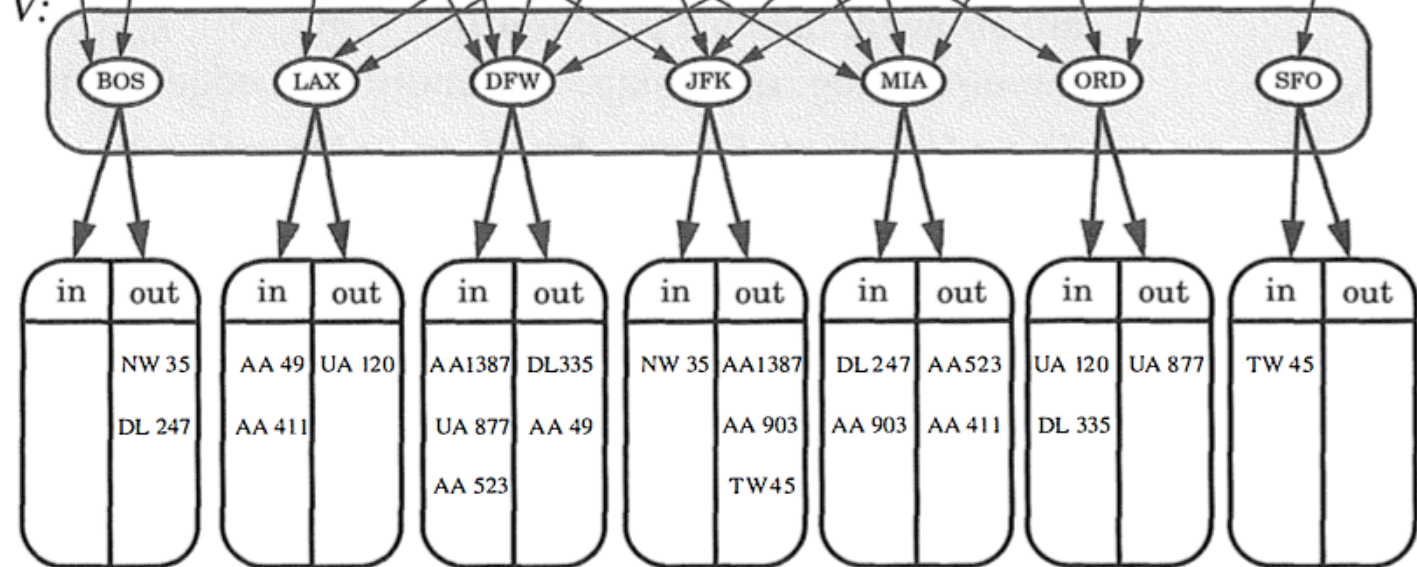
A container of edge objects, where each edge object references the origin and destination vertex object

# Adjacency List Structure

*E*:



*V*:



An edge list structure, where additionally each vertex object  $v$  references an incidence container which stores references to the edges incident on  $v$ .

# Adjacency Matrix Structure

		0	1	2	3	4	5	6
		BOS	DFW	JFK	LAX	MIA	ORD	SFO
		0	1	2	3	4	5	6
BOS	0	∅	∅	NW 35	∅	DL 247	∅	∅
DFW	1	∅	∅	∅	AA 49	∅	DL 335	∅
JFK	2	∅	AA 1387	∅	∅	AA 903	∅	TW 45
LAX	3	∅	∅	∅	∅	∅	UA 120	∅
MIA	4	∅	AA 523	∅	AA 411	∅	∅	∅
ORD	5	∅	UA 877	∅	∅	∅	∅	∅
SFO	6	∅	∅	∅	∅	∅	∅	∅

A 2D array of all vertex pairs, where cell  $A[u,v]$  stores edge  $e$  incident on vertices  $u,v$  if such an edge exists.

# Asymptotic Performance

<ul style="list-style-type: none"> <li>◆ <math>n</math> vertices, <math>m</math> edges</li> <li>◆ no parallel edges</li> <li>◆ no self-loops</li> <li>◆ Bounds are “big-Oh”</li> </ul>	Edge List	Adjacency List	Adjacency Matrix
Space	$n + m$	$n + m$	$n^2$
incidentEdges( $v$ )	$m$	deg( $v$ )	$n$
areAdjacent ( $v, w$ )	$m$	min(deg( $v$ ), deg( $w$ ))	1
insertVertex( $o$ )	1	1	$n^2$
insertEdge( $v, w, o$ )	1	1	1
removeVertex( $v$ )	$m$	deg( $v$ )	$n^2$
removeEdge( $e$ )	1	1	1